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An energy-efficient gated Johnson counter using redundant switching technique

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Abstract. This paper proposes an approach to increase the energy efficiency of gated Johnson counters by eliminating the redundant switching consumption of the unused counter-chain cells. Two techniques that aim to achieve this goal are analysed and show that efficiency improves almost exponentially as the counter resolution increases up to a limiting asymptote. To demonstrate the effectiveness of the proposed approach, a 5-bit Johnson counter architecture with 16 stages is designed. The design is extensively analysed and simulated up to 1 GHz clock frequency and implemented into single channel SS-ADC as part of CIS realized in a 0.11 μ m 1P4M BSI CMOS process. The obtained results show that the proposed approach trades energy efficiency for the area and is practically suitable only for counters with resolution ≤ 8 bits.

1. Introduction

As the size of systems-on-chip continues to shrink and system speeds increases, low-power circuit design has become a critical concern in VLSI design. This is especially true for synchronous sequential circuit logic, which is widely used in various systems, including image sensors, and areas where micro-electronics reliability is paramount importance [1]. Among the conventional building blocks from sequential logic, counters are a core component, and their power efficiency has a huge impact on system performance [2], [3].

With the increase of electronic devices in the consumer basket the transfer from standalone electronics into integrated requires more digital knowledge and experience into microelectronics design furthermore in sequential logic in order to meet labor market needs in field of microelectronic industry in which Johnson counters are heavily represented, are purposefully included, and studied in the university curriculum [4], [5].

Some conventional counter architectures are typically designed using either binary or thermometer output coding in a synchronous or asynchronous switching domain. However, power inefficiencies can arise from unnecessary gate activity, leading to high power dissipation. Most efforts for clock power management in systems are focused on reducing voltage swings, optimal buffering and clock distribution and routing [6]. However, oftentimes, inefficiency comes from unnecessary gate activity from where the clock gating techniques as such have been extensively analyzed by researchers [7][8].

The aim of this paper is to propose a low-power technique applied to a Johnson counter that reduces the size of the whole counter chain and increases consumption efficiency without speed degradation. The studied system has lower power dissipation and a simpler design compared to the conventional one. The verification of the system efficiency is achieved by designing and comparing various counter architectures, followed by additional power analysis verified with SPICE engine and

further implemented into single channel single slope ADC architecture that takes place into final imager CIS product.

2. Architecture overview

The proposed Johnson counter architecture is displayed in Fig. 1 that is also known as a walking-ring counter that output is achieved by a thermometer code structure. The counter implementation consists of a chain of serially connected modified D flip-flop (DFF) cells.

In Fig. 2 reveals a) a single cell from a Johnson counter that is composed of a master-slave DFF based on tri-states structure with included extra combinatorial logic in the periphery. The modified DFF cell includes in-self OR, NAND and MUX gates that components forming an additional combinatorial logic. The logic that is formed by adding extra logic leads to achieve gating technique that generates a local clock pulses for each modified cells, that further leads eliminates redundant switches in the counter chain. Implementing such a technique results in increased consumption efficiency with an increase in resolution, making the counter consumption approximately fixed and not strongly resolution-dependent.

In order to enable clock gating feature to be compatible with double data rate (DDR) counting, the periphery logic is design in such a way that two clock paths is performed to a clock input of conventional DFF. Each clock path enable is self-configured that is strongly dependent on the input data so that each periphery cell receives feedback from the inverted output of the DFF, the input data (from previous cell), and the state of the main (master) clock. The inverter shown in the block diagram is not a physical component but shows the logical inversion of a clock that is sensed to a NAND input. The inversion is used only in one of the inputs of NAND path in order to maintain the same clock phase (polarity) of both paths and no glitches occur by switching between both clock paths.

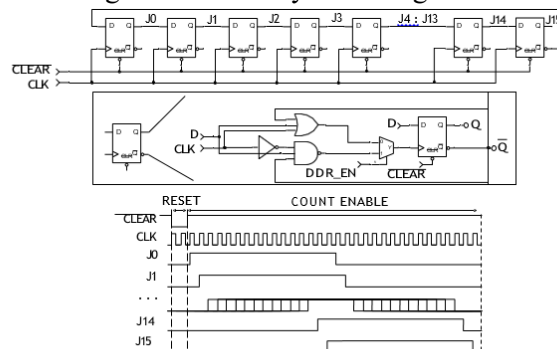
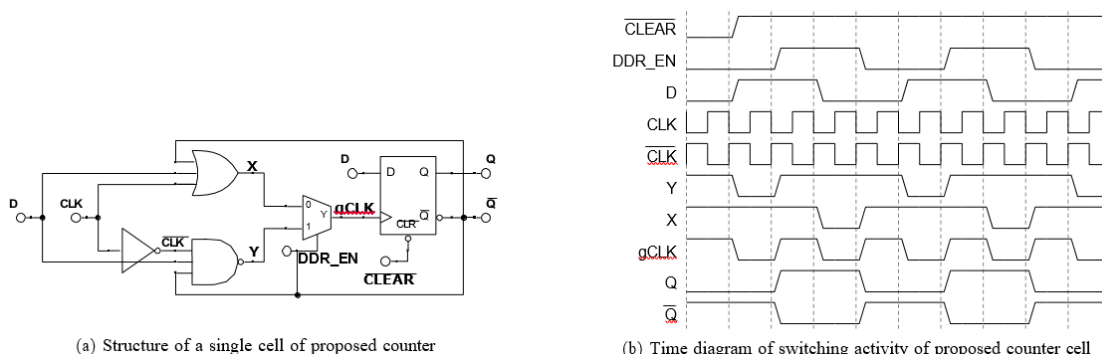


Fig. 1: Block diagram of a Johnson counter



(a) Structure of a single cell of proposed counter

(b) Time diagram of switching activity of proposed counter cell

The timing diagram is displayed in Fig. 2 b) that reflects the operation principle of a single channel (cell) gated counter. The individual cell is initialized by toggling to a low state of a *CLEAR* signal so that the output is referred respectively to $Q = 0$ and $Q = 1$. The initial state is followed by switching activity that fills the counter with 'ones' where the local clock is being generated by the NAND gate noted by *Y*. The switching activity continues until counter overflow occur that leads to toggle last cell in the DFF chain to be inverted so that output Q that is

shorted to the first cell input realizes DDR counting. The following counting activity is replaced by feeding 'zeroes' in the chain. Although the local clock is now gated by *OR* instead of *NAND* noted by *X*. The generated *gCLK* is sensed to the clock input of DFF by *MUX* whose control signal is self-triggered by *Q* of the last chain cell. The intuitively proposed operation is extensively analyzed and described in following section in order to verify efficiency of the architecture.

3. Energy Analysis

To effectively analyze the system, it is recommended to split the process into firstly analyzing the DFF displayed in Fig. 3 as an individual cell followed by analyze with including those extra combinatorial logic. The structure of the DFF used in proposed architecture is composed of tri-state inverters, conventional CMOS inverters, NAND, and complementary switches. The consumption analysis of the DFF is separated into two main aspects based on:

- Static activity of DFF - this occurs when the previous and present DFF are in same state. The power consumption is only determined by the master sub-latch that contains two face-to-face tri-state inverters and conventional CMO inverter. When input data of the DFF is unchanged to slave keeps the it's storage sot that no switching activity occur hence no consumption ($P_{Slave} = 0W$). That statement can be concluded as:

$$P_{nonSw} = P_{Master} + P_{Slave} = 2P_{tri-state} + P_{inv} \quad (1)$$

where P_{nonSw} , P_{Master} , P_{Slave} represents non-switching power consumption that is the sum of master and slave latch consumption.

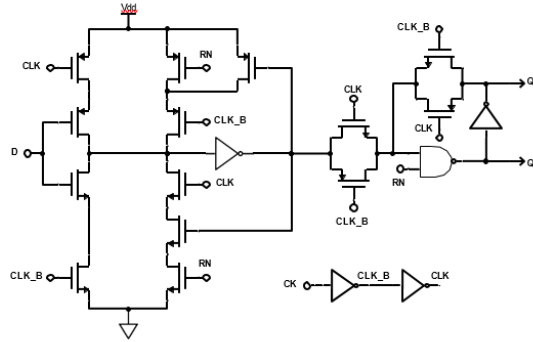


Fig. 3: Structure of DFF

With including clock gating technique with extra combinatorial periphery the redundant switching of cells where data is remain unchanged is excluded so that the consumption of the clock is determined only by DFF that dynamically switching.

- Dynamic activity of DFF - Also known as switching activity that occurs when the present data is opposite to the previous data and vice versa. That activity causes both DFF halves (master and slave) to consume noted as dynamic (switching or toggling) power. Such a configuration leads increase maximum consumption that can be concluded as:

$$P_{Sw} = P_{Master} + P_{Slave} = 2P_{tri-state} + 2P_{inv} + P_{NAND} \quad (2)$$

In a free running Johnson counter is mostly composed of single cell that is dynamically switching by nature of the architecture. Thus helps to keep the power consumption to be resolution independent. This statement is true up-to 8-bit resolution. With increasing chain the number of cells drastically increases that leads to beefy drivers to drive the global signal to the counter cells (that most of consumption comes from).

Further consumption analyses is needed with adding extra combinatorial logic in order to determine the overall consumption of single cell that is equal to overall consumption of the counter

system. The consumption of individually switching cell is determined by digital component that is along the clock path that also depends on input feeding 'ones' or 'zeroes' that is concluded as follows:

- During the initial state of operation, the all cells in the chains are in reset, causing the output of the DFFs to be $Q = 0$ and $Q = 1$. The operation of the counter starts with feeding 'ones' and performing clock path through NAND gate. So that switching consumption is determined by:

$$P_{ones} = P_{NAND} + P_{DFF} \quad (3)$$

where P_{ones} , P_{NAND} and P_{DFF} are respectively consumption of the NAND and DFF. That operation continues until all cells are overflowed by 'ones' so that second stage of operation occur.

- The second stage operation of the counter starts with feeding 'zeroes' and performing clock path through OR gate. So that switching consumption is determined by:

$$P_{zeroes} = P_{OR} + P_{DFF} \quad (4)$$

where P_{zeroes} , P_{OR} and P_{DFF} are respectively consumption of the OR and DFF. That operation stage continues until all cells are under-flowed by 'zeroes' so that counter starts from beginning.

To sum up the overall power consumption P_{prop} is equal to:

$$P_{prop} = n(P_{ones} + P_{zeroes}) \quad (5)$$

where n is counter resolution.

To further increase the power consumption using of a DDR technique is a must that leads to reduction number of cells in the chain by half so that resolution is reduced to $2n-1$ from $2n$ that also leads to compact layout design.

4. Experimental results

To assess the counter efficiency of the proposed techniques compared to conventional ripple-carry (asynchronous), synchronous binary counter, conventional synchronous Johnson counter is shown in the timing diagram in Fig.4 where the current magnitude during a transition toggle event between

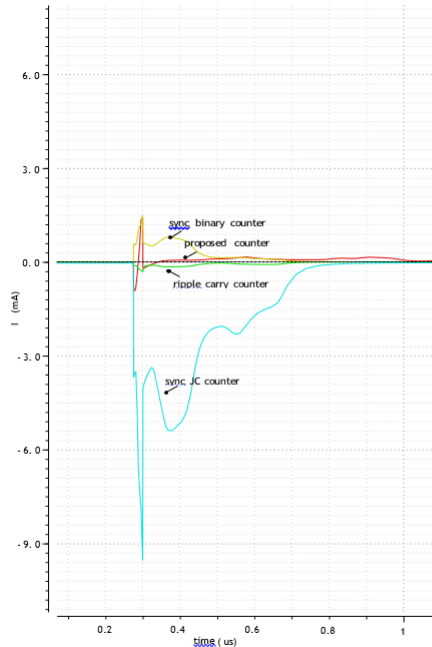
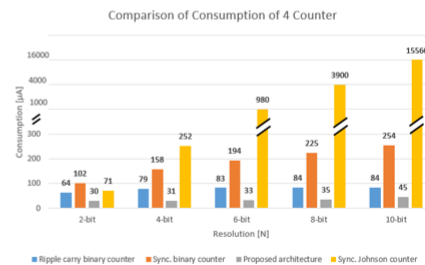


Fig. 4: Comparison between counter architectures: a) Ripple carry counter; b) Synchronous binary counter; c) Conventional Johnson counter; d) Proposed counter

Comparison Table						
Resolution	Name of counter	# of DFF	Consumption [µA]	PD (from : to)	size [µm]	
2-bit	Ripple carry binary counter	2	64	150 - 350	70	
	Sync. binary counter	2	102	170	70	
	Sync. Johnson counter	3	71	170	105	
	Proposed architecture	2	30	370	70	
4-bit	Ripple carry binary counter	4	79	150 - 700	140	
	Sync. binary counter	4	158	170	140	
	Sync. Johnson counter	15	252	170	525	
	Proposed architecture	8	31	370	210	
6-bit	Ripple carry binary counter	6	83	150 - 1000	210	
	Sync. binary counter	6	194	170	210	
	Sync. Johnson counter	63	980	170	2205	
	Proposed architecture	32	33	370	1120	
8-bit	Ripple carry binary counter	8	84	150 - 1400	280	
	Sync. binary counter	8	225	170	280	
	Sync. Johnson counter	255	3900	170	8925	
	Proposed architecture	128	35	370	4480	
10-bit	Ripple carry binary counter	10	84	150 - 1800	350	
	Sync. binary counter	10	254	170	350	
	Sync. Johnson counter	1023	15560	170	35805	
	Proposed architecture	512	45	370	17920	

(a) Comparison Table Results



(b) Comparison Graph of Consumption of 4 Counters

Fig. 5: Counter Comparison Table and Graph

”01111” and ”10000” of all examined counter types. It is clearly evident that the synchronous binary performs worst, while the Johnson clock gated counter offers the best power efficiency.

Due to the gating method, the consumption efficiency increases with an increase in the counter resolution. To verify this statement 2, 4, 6, 8, and 10-bit counters architectures were designed and simulated. The results are shown in table in Fig. 5 a) and a consumption comparison graph in Fig 5 b). The graphs as a function of resolution indicate that power consumption with the clock gating method is kept relatively constant with an increase in resolution. All results were obtained from simulations using the Spectre tool with a supply voltage of 1.5V, typical 1P4M 110nm CMOS process, at room temperature, with typical process parameters with a capacitive load of 50fF at output count nodes.

The proposed 5-bit clock-gated double data rate counter is physically designed with overall sizes $40\mu\text{m} \times 55\mu\text{m}$. The compact single cell size is $18.5\mu\text{m} \times 4\mu\text{m}$.

5. Conclusion

A vastly important conclusion to draw from the presented analysis is that the clock gating method applied to Johnson counting, although improving greatly power consumption, comes at a significant cost of silicon area compared to binary counting. This is probably the largest trade-off of this technique, which makes it limiting in practical applications that require counter resolutions greater than 8-bits.

Regardless of the above, this work proposed a CMOS synchronous counting method that reduces redundant switching power consumption during counting, resulting in energy efficiency. The conducted simulations and system analysis, as well as comparison between well-known four different 5-bit architectures with 1 GHz count speed verifies the design.

The whole system can be a standalone counter or it can take place in a larger system such as time or analog-to-digital converter, sequencer or even in CPU. Proposed counter architecture is implemented as a part of a single slope analog-to-digital converter.

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